UPPER BOUNDS FOR INTEGER SOLUTIONS TO A SYSTEM OF TWO BILINEAR FORMS

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ABSTRACT. We show that the number of integer solutions for a pair of bilinear equations in at least 2×6 variables has (up to logarithms) the expected upper bound unless there is a structural reason why it is not the case.

1. Introduction

In the by now classical work, Birch [1] provides a method to show that a system of forms of degree d has the expected number of solutions as long as the number of variables is big enough compared to the dimension of the 'singular locus'. If we have a system of R forms in n variables and V^* is the 'singular locus', then the condition is

$$n - dim[V^*] > R(R+1)(d-1)2^{d-1}$$
.

In this paper we want to consider the case of two bilinear forms in 2s variables. The condition in this case would be $2s - dim[V^*] > 12$. It can probably be improved by recent work of Schindler [3] on Birch's theorem for bihomogeneous forms.

From a naïve point of view, it seems very strange that the method should give weaker results if we find ourselves in more structured situations with a large singular locus, like the case of two diagonal forms, where the singular locus is at least as big as s. On the other hand, the standard circle method approach is well suited to answer the diagonal problem in as few as 2×4 variables.

The question we are trying to answer is: Can we prove a result for 'all' bilinear forms, independent of the size of the singular locus? We show that this is indeed the case if we restrict our attention to upper bounds instead of asymptotic formulas.

Theorem 1.1. Let $N \geq 2$ and $B_1, B_2 \in \mathbb{Z}^{s \times s}$ be two integer matrices with $B_1(\mathbf{x}, \mathbf{y}) = \mathbf{x}^T B_1 \mathbf{y}$, $B_2(\mathbf{x}, \mathbf{y}) = \mathbf{x}^T B_2 \mathbf{y}$ corresponding bilinear forms, then the number of solutions to the system $B_1(\mathbf{x}, \mathbf{y}) = B_2(\mathbf{x}, \mathbf{y}) = 0$ with $|x_i| \leq N, |y_i| \leq N$ is of order $O(N^{2s-4}(\log N)^2)$ as long as we are not in one of the two situations

- (1) $rank(\lambda B_1 + \mu B_2) \le 5$ for all $\lambda, \mu \in \mathbb{Z}$ or
- (2) $rank(\lambda B_1 + \mu B_2) \leq 1$ for some $(\lambda, \mu) \neq (0, 0)$.

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Remark 1.2. The number of solutions is bounded from below by cN^{2s-4} for some c > 0 by an averaging argument. See also Lemma 2.5 below.

The first exceptional case in Theorem 1.1 does not give a sharp theoretical bound, but reflects a limitation in our methods. We would expect that the same result holds with a three replacing the five (which would be best possible).

The second roughly corresponds to one of the equations being of the form xy = 0, which cannot 'save' two variables, as required in the theorem (see also Lemma 2.5). In this case our result is best possible.

Rank conditions as those that appear in our theorem are typical in this line of work as can be seen in previous work of Schmidt [4] and Dietmann [2], who deal with systems of general quadratic forms.

Most of the following arguments will extend to general systems of bilinear forms, but we feel that the methods and ideas are best presented in the simplest case of two equations.

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2. Collecting the Tools

Before we begin with stating the main lemmata of this work, we need a few notational conventions.

As usual, we use O-notation and the Vinogradov notation $f \ll g$ to denote that $|f| \leq C|g|$ for some C > 0. In the same way, we say that the number of solutions S(N) is essentially bounded by a quantity T(N), if there is a C > 0 such that $S(N) \leq T(CN)$ is a bound for all $N \in \mathbb{N}$.

Now we want to state the tools that we are going to use excessively throughout the paper. Most of them are simple results from linear algebra.

Lemma 2.1 (Homogenisation). Let $A \in \mathbb{Z}^{s \times s}$, $\mathbf{c} \in \mathbb{Z}^s$ and $y_i \in \mathbb{Z}$, $|y_i| \leq N$.

- (i) Let $A\mathbf{y} = \mathbf{c}$ be a system of inhomogeneous linear equations. Then the number of solutions to this equation is essentially bounded by the number of solutions to the homogeneous system $A\mathbf{y} = \mathbf{0}$.
- (ii) Let $A\mathbf{y} = \mathbf{0}$ be a system of linear equations. Then the number of solutions to this equation is essentially bounded by N times the number of solutions to the same system with $y_j = 0$ for some $1 \le j \le s$.
- (iii) If the last d entries of $A\mathbf{y}$ don't depend on the variables y_1, \ldots, y_j , then we can set the variables y_{j+1}, \ldots, y_s equal to zero in the first r-d equations of $A\mathbf{y}=\mathbf{0}$ and obtain an essential upper bound for the number of solutions. In other words: If A is a upper triangular block matrix, we can change it into a diagonal block matrix.

Remark 2.2. A bilinear system $\mathbf{x}^T B_i \mathbf{y} = 0$ can always be thought of as a linear system in \mathbf{y} by fixing the variables \mathbf{x} (or the other way around).

Proof. For the first statement, we observe that for a given fixed solution $A\mathbf{z} = \mathbf{c}$ and any other solution $A\mathbf{y} = \mathbf{c}$ to the inhomogeneous linear equation, we obtain a solution $A(\mathbf{y} - \mathbf{z}) = \mathbf{0}$ with $\|\mathbf{y} - \mathbf{z}\|_{\infty} \leq 2N$.

For the second statement, we observe that by fixing y_j , we can rewrite $A\mathbf{y} = \mathbf{0}$ into $B\mathbf{y}' = \mathbf{c}_j$, where B is essentially A but with missing column j and \mathbf{c}_j is $-y_j$ times the jth column of A. The result follows from part one and the observation, that there are O(N) choices for y_j .

The third statement is slightly more difficult. For any choice of values for $\mathbf{y}'' = (y_{j+1}, \ldots, y_s)$ that satisfy the last d equations of $A\mathbf{y} = \mathbf{0}$, we can set $\mathbf{y}' = (y_1, \ldots, y_j)$ and write the first r - d equations in the form $B\mathbf{y}' = c(\mathbf{y}'')$, where B is the upper left submatrix of size $(r - d) \times j$. By part (i), this system is majorized by the system $B\mathbf{y}' = 0$. This homogenisation procedure doesn't affect the last d equations since they are independent of \mathbf{y}' .

Lemma 2.3 (Divisor estimates). An equation of the form $dx_1y_1 = cx_2y_2$ with $c, d \neq 0$ has $O(N^2 \log N)$ solutions with $|x_i|, |y_i| \leq N$.

Proof. A Cauchy-Schwarz symmetrisation (see proof of Lemma 2.6 below) reduces the problem to the form $x_1y_1 = x_2y_2$. For $x_1 = 0$, the number of solutions is $O(N^2)$. We can therefore assume that x_1 and x_2 are nonzero and positive. If we consider this to be a linear equation in y_1 and y_2 and set $d = \gcd(x_1, x_2)$, then we can instead look at $u_1y_1 = u_2y_2$, where $\gcd(u_1, u_2) = 1$ and $u_i = x_i/d$. This forces the divisibility conditions $u_1|y_2$ and $u_2|y_1$. Therefore, the number of solutions to this linear diophantine equation is bounded by $(2N+1)/\max(u_1, u_2)$. We obtain an essential upper bound of the form

$$N \sum_{1 \le x_1, x_2 \le N} \frac{\gcd(x_1, x_2)}{\max(x_1, x_2)}.$$

Collecting the terms with equal greatest common divisor, we obtain

$$N \sum_{d \le N} \sum_{1 \le u_1, u_2 \le N/d} \frac{1}{\max(u_1, u_2)} \le 2N \sum_{d \le N} \sum_{1 \le u_1 \le u_2 \le N/d} \frac{1}{u_2}$$
$$=2N \sum_{d \le N} \sum_{1 \le u_2 \le N/d} 1 = 2N \sum_{d \le N} \frac{N}{d} \ll N^2 \log N.$$

The next lemma isn't strictly necessary for the argument, but simplifies the exposition.

Lemma 2.4 (Coordinate change). If we set $\mathbf{u} = K\mathbf{x}$ for a matrix $K \in \mathbb{Z}^{s \times s}$ of full rank, then the resulting system $\mathbf{u}K^{-1}B\mathbf{y} = \mathbf{0}$ has a bigger upper bound, as long as we choose C > 0 with $|u_i| \leq CN$ in such a way that $[-CN, CN]^s$ covers the image of [-N, N] by K. By multiplication with suitable integers, we can also assume that the coefficients of the new system are integers.

Proof. Every solution in **x** translates into a solution in **u**.

Lemma 2.5 (Lower bound). A system $\mathbf{x}^T B_j \mathbf{y} = 0$ with $1 \leq j \leq r$ in 2s variables has $\gg N^{2(s-r)}$ many solutions with $|x_i|, |y_i| \leq N$.

Proof. Lemma 2.1 shows us that the number of solutions to the system $\mathbf{x}^T B_j \mathbf{y} = h_j$ for arbitrary fixed h_j is essentially bounded by the number of solutions to the system $\mathbf{x}^T B_j \mathbf{y} = 0$. If we now consider h_j to be variables as well, which have the range $|h_j| \leq C_j N^2$ for some large enough $C_j > 0$ (depending on B_j), then the total number of solutions to the system $\mathbf{x}^T B_j \mathbf{y} = h_j$ is N^{2s} since we can choose x_i and y_i freely and this choice fixes the values of all h_j . We obtain

$$N^{2s} = \sum_{\mathbf{h}} \#\{x_i, y_i : \mathbf{x}^T B_j \mathbf{y} = h_j\} \ll \sum_{\mathbf{h}} \#\{x_i, y_i : \mathbf{x}^T B_j \mathbf{y} = 0\}$$

$$\ll N^{2r} \#\{x_i, y_i : \mathbf{x}^T B_j \mathbf{y} = 0\}.$$

Lemma 2.6 (Diagonal system). The system

$$d_1x_1y_1 + d_2x_2y_2 = d_3x_3y_3 + d_4x_4y_4,$$

$$e_1x_1y_1 + e_2x_2y_2 = e_3x_3y_3 + e_4x_4y_4,$$

has $O(N^4(\log N)^2)$ solutions with $|x_i|, |y_i| \le N$ if and only if every 2×3 submatrix of

$$\begin{pmatrix} d_1 & d_2 & d_3 & d_4 \\ e_1 & e_2 & e_3 & e_4 \end{pmatrix}$$

has rank two.

Proof. Let us first assume that every 2×3 submatrix has rank two. This implies that we can rearrange matters such that the submatrices $\begin{pmatrix} d_1 & d_2 \\ e_1 & e_2 \end{pmatrix}$

and $\begin{pmatrix} d_3 & d_4 \\ e_3 & e_4 \end{pmatrix}$ have rank two. Write the number of solutions to the system as a sum and perform a simple Cauchy-Schwarz symmetrisation of the coefficients.

$$\sum_{|\mathbf{x}| \le N} \sum_{|\mathbf{y}| \le N} \sum_{\substack{d_1 x_1 y_1 + d_2 x_2 y_2 = n = d_3 x_3 y_3 + d_4 x_4 y_4 \\ e_1 x_1 y_1 + e_2 x_2 y_2 = m = e_3 x_3 y_3 + e_4 x_4 y_4}} 1$$

$$= \sum_{n,m} \left(\sum_{\substack{d_1 x_1 y_1 + d_2 x_2 y_2 = n \\ e_1 x_1 y_1 + e_2 x_2 y_2 = m}} 1 \right) \cdot \left(\sum_{\substack{d_3 x_3 y_3 + d_4 x_4 y_4 = n \\ e_3 x_3 y_3 + e_4 x_4 y_4 = m}} 1 \right)$$

$$\le \left(\sum_{n,m} \left(\sum_{\substack{d_1 x_1 y_1 + d_2 x_2 y_2 = n \\ e_1 x_1 y_1 + e_2 x_2 y_2 = m}} 1 \right)^2 \right)^{1/2} \cdot \left(\sum_{n,m} \left(\sum_{\substack{d_3 x_3 y_3 + d_4 x_4 y_4 = n \\ e_3 x_3 y_3 + e_4 x_4 y_4 = m}} 1 \right)^2 \right)^{1/2}$$

The first term corresponds to the system

$$d_1x_1y_1 + d_2x_2y_2 = d_1x_3y_3 + d_2x_4y_4,$$

$$e_1x_1y_1 + e_2x_2y_2 = e_1x_3y_3 + e_2x_4y_4,$$

and by taking linear combinations, we can reduce this to

$$x_1 y_1 = x_3 y_3, x_2 y_2 = x_4 y_4,$$

which has the required number of solutions by Lemma 2.3. An equivalent argument takes care of the other term.

In the case that there is a 2×3 submatrix with rank at most one, we can simplify the system to

$$d_1x_1y_1 + d_2x_2y_2 = d_3x_3y_3 + d_4x_4y_4,$$

$$0 = e_4x_4y_4.$$

If $e_4 = 0$ we can show by the argument from Lemma 2.5 that the first equation has $\gg N^6$ solutions. In the case $e_4 \neq 0$ we have $x_4 = 0$ (or $y_4 = 0$), which reduces the first equation to six variables and a factor of $\gg N$ from the summation over y_4 (or x_4). Again by Lemma 2.5 we have at least $\gg N^5$ solutions.

The last ingredient is another simple observation about systems of linear equations.

Lemma 2.7. Let $A\mathbf{y} = \mathbf{c}$ be a linear equations system, where all equations are independent of y_s apart from the last equation. In other words, the last column of A is a non-zero multiple of the standard basis vector \mathbf{e}_s . Then the number of solutions to this system is bounded by the number of solutions to $A'\mathbf{y}' = \mathbf{c}'$, where we get A' by removing the last column of A and \mathbf{y}', \mathbf{c}' by removing the last entry of \mathbf{y} and \mathbf{c} .

Proof. For fixed values of the variables y_1, \ldots, y_{s-1} , there is at most one value of y_s that makes the last equation true.

3. The starting point

It turns out that it is sufficient to prove the main result for s=6 in order to get it for all $s \ge 6$ as we will see in Section 7. We therefore begin with the system of two bilinear equations in two times six variables

$$B_1(\mathbf{x}, \mathbf{y}) = B_2(\mathbf{x}, \mathbf{y}) = 0.$$

By taking linear combinations of the two equations, we can assume that $rank(B_1) = 6$ as long as we are not in the first exceptional case of Theorem 1.1.

One way to look at the system is to consider them as linear equations in \mathbf{y} with coefficients being linear forms in \mathbf{x} . We get

$$K_1(\mathbf{x})y_1 + K_2(\mathbf{x})y_2 + K_3(\mathbf{x})y_3 + K_4(\mathbf{x})y_4 + K_5(\mathbf{x})y_5 + K_6(\mathbf{x})y_6 = 0,$$

 $L_1(\mathbf{x})y_1 + L_2(\mathbf{x})y_2 + L_3(\mathbf{x})y_3 + L_4(\mathbf{x})y_4 + L_5(\mathbf{x})y_5 + L_6(\mathbf{x})y_6 = 0.$

By a change of coordinates (Lemma 2.4), we can assume that the first equation is diagonal. This simplifies the situation to

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$L_1(\mathbf{x})y_1 + L_2(\mathbf{x})y_2 + \ldots + L_5(\mathbf{x})y_5 + L_6(\mathbf{x})y_6 = 0.$$

The main difficulty to overcome is the interdependence of the two equations. Our goal will be to either extract a system with separated variables or one with a diagonal structure.

One possible way to force independence is to set set $x_1 = x_2 = 0$. With $\mathbf{x}' = (0, 0, x_3, x_4, x_5, x_6)$ Lemma 2.1 (ii) gives us a factor of $O(N^2)$ and the system

$$x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$L_1(\mathbf{x}')y_1 + L_2(\mathbf{x}')y_2 + L_3(\mathbf{x}')y_3 + \dots + L_6(\mathbf{x}')y_6 = 0.$$

This makes the first equation independent of y_1 and y_2 . By Lemma 2.1 (iii), we can remove the dependence of the second equation on y_3, \ldots, y_6 , which leaves us with

$$x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

 $L_1(\mathbf{x}')y_1 + L_2(\mathbf{x}')y_2 = 0.$

Now that the second equation is independent of y_5 and y_6 , we can do the same thing to the first line and obtain the majorising system

$$x_5 y_5 + x_6 y_6 = 0,$$

$$L_1(\mathbf{x}') y_1 + L_2(\mathbf{x}') y_2 = 0.$$

If L_1 and L_2 depend on (x_3, x_4) in a non-singular way (see below), a final change of variables would give us the system

$$x_5y_5 + x_6y_6 = 0,$$

$$x_3y_1 + x_4y_2 = 0.$$

Since we achieved independence, Lemma 2.3 gives us $O(N^4(\log N)^2)$ solutions. To obtain the final bound, we collect the $O(N^2)$ contribution from our first step and the $O(N^2)$ from the sum over (y_3, y_4) .

Let us explore the conditions under which the above argument works. Write l_{ij} for the jth coefficient of L_i . So that $L_1(\mathbf{x}) = l_{11}x_1 + \ldots + l_{16}x_6$ and $L_2(\mathbf{x}) = l_{21}x_1 + \ldots + l_{26}x_6$. If the matrix $\begin{pmatrix} l_{13} & l_{14} \\ l_{23} & l_{24} \end{pmatrix}$ has rank two, then the change of variables from $(L_1(\mathbf{x}'), L_2(\mathbf{x}'))$ to (x_3, x_4) will be successful.

To understand the complementary case, we observe that we made some arbitrary choices along the way. Consider the 'off-diagonal' matrix build from the coefficients of the linear forms L_1 and L_2 , given by

$$\begin{pmatrix} l_{13} & l_{14} & l_{15} & l_{16} \\ l_{23} & l_{24} & l_{25} & l_{26} \end{pmatrix}.$$

Whenever this matrix has full rank, the above strategy will also work by choosing a a possibly different pair of indices than 3 and 4, which corresponds to the special case that the first two columns are linearly independent. So in order for this not to work, we need that matrix (3.1) has rank at most one.

On the other hand, we can set any pair of variables $\{x_i, x_j\}$ equal to zero in the first step of the argument, not necessarily x_1 and x_2 . Since the matrix (3.1) sits in the upper right corner of B_2 , this translates (by permuting the variables) into the following rank condition for the matrix B_2 : Any off-diagonal matrix in B_2 has rank at most one. ('Off-diagonal' means that it doesn't contain any diagonal elements.)

Write $\mathbf{v} \otimes \mathbf{w} := \mathbf{v} \cdot \mathbf{w}^T$, \mathbf{e}_i to be the *i*th standard basis vector and define the *off-rank* of a matrix to be the maximal rank of an off-diagonal submatrix. We have the following classification of off-rank one matrices.

Lemma 3.1. A matrix $B \in \mathbb{Z}^{s \times s}$ with off-rank one has the form

- (i) $B = D + \mathbf{v} \otimes \mathbf{w}$,
- (ii) $B = D + \mathbf{v} \otimes \mathbf{e}_i + \mathbf{e}_i \otimes \mathbf{w}$ or
- (iii) B = D + E,

where D is a diagonal matrix, $\mathbf{v}, \mathbf{w} \in \mathbb{Q}^s$ and $E \in \mathbb{Z}^{s \times s}$ has non-zero entries only in a 3×3 submatrix, which is based on the diagonal.

Proof. See Appendix A.
$$\Box$$

Example 3.2. To get a better feeling for this concept, we give examples of the three possible cases.

$$\begin{pmatrix} 0 & 2 & 2 & 2 & 1 \\ 3 & -1 & 6 & 6 & 3 \\ 0 & 0 & 6 & 0 & 0 \\ 1 & 2 & 2 & -5 & 1 \\ 1 & 2 & 2 & 2 & 3 \end{pmatrix}, \begin{pmatrix} 3 & 10 & 0 & 0 & 0 \\ 3 & 7 & 5 & 2 & 8 \\ 0 & -2 & 5 & 0 & 0 \\ 0 & -3 & 0 & 1 & 0 \\ 0 & 2 & 0 & 0 & -4 \end{pmatrix}, \begin{pmatrix} 1 & 17 & 2 & 0 & 0 \\ 3 & -1 & 1 & 0 & 0 \\ 4 & -1 & 6 & 0 & 0 \\ 0 & 0 & 0 & 6 & 0 \\ 0 & 0 & 0 & 0 & 2 \end{pmatrix}.$$

We have $\mathbf{v} = (1, 3, 0, 1, 1)^T$ and $\mathbf{w} = (1, 2, 2, 2, 1)^T$ for the first example.

Each of the next three sections is dealing with one of the cases in 3.1.

Remark 3.3. The off-rank zero case for B_2 is covered by any of the following sections.

In order to understand the structure of $B_2 = D + \mathbf{v} \otimes \mathbf{w}$ we introduce new variables $h = \mathbf{x}^T \mathbf{v}$ and $l = \mathbf{w}^T \mathbf{y}$. Then the bilinear system transforms into

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 + d_2x_2y_2 + d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = hl,$$

$$v_1x_1 + v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_2y_2 + w_3y_3 + w_4y_4 + w_5y_5 + w_6y_6 = l.$$

This system has now the advantage of being diagonal, while having a higher complexity due to the two additional linear equations.

The exact behaviour of this system depends on the coefficients d_i , v_i and w_i . We use Lemma 2.1 to set $x_1 = y_2 = 0$ similar to the procedure in the non-degenerate case in Section 3. The linear equations in the system

$$x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = hl,$$

$$v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_3y_3 + w_4y_4 + w_5y_5 + w_6y_6 = l,$$

can be dealt with by Lemma 2.7 as long as the coefficients v_2 and w_1 are non zero. We end up with the reduced problem of bounding the solutions to

$$x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = hl.$$

By Lemma 2.6 (and Lemma 2.1) we have $O(N^6(\log N)^2)$ solutions as long as not all d_i are equal. Together with the $O(N^2)$ contribution from Lemma 2.1 in the first step, we obtain the result.

As in the previous section, we need to analyse the argument to obtain a good description of the complementary case. The method words if $v_2, w_1 \neq 0$ and $d_i \neq d_j$ for some $i, j \in \{3, 4, 5, 6\}$. By symmetry (remaining of variables), we can perform the argument with different sets of indices as well.

The first step succeeds, therefore, if there are v_i and w_j with $i \neq j$, which are both non-zero. Let us explore the complementary situation.

Case 1: $v_i \cdot w_j = 0$ for all $i \neq j$.

This implies that either $\mathbf{v} = \mathbf{0}$, $\mathbf{w} = \mathbf{0}$ or that \mathbf{v} and \mathbf{w} have only one non-zero component with the same index.

Case 1.1: $\mathbf{w} = \mathbf{0}$ ($\mathbf{v} = \mathbf{0}$ is equivalent by symmetry). Since $B_2 = D + \mathbf{v} \otimes \mathbf{w}$, we get the system

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 + d_2x_2y_2 + d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = 0.$$

Lemma 2.6 and Lemma 2.1 give the right answer as long as the d_i take on three different values. If there are only two different values for d_i , at least three of the coefficients have to be the same and a linear combination (with a renaming of variables) brings us to

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = 0,$$

where $d_4, d_5 \in \{0, d_6\}$. By Lemma 2.1 (iii) this simplifies further to

$$x_1y_1 + x_2y_2 = 0 = d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6.$$

If $d_5 = d_6 \neq 0$ or $d_4 = d_6 \neq 0$, then we get the correct upper bound by Lemma 2.3. Otherwise, we end up with at most one non-zero coefficient, which brings us into the second exceptional case of Theorem 1.1.

Case 1.2: $w_j = v_i = 0$ for $i, j \ge 2$ (similar cases by coordinate change). The linear equations simplify to $v_1x_1 = h$ and $w_1y_1 = l$ and the whole system changes into

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$(d_1 - v_1w_1)x_1y_1 + d_2x_2y_2 + d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = 0.$$

This is the same situation we faced in the previous case and can be dealt with accordingly.

Now we are going to discuss the second part of the general argument in this section, where we needed that at least one of the coefficients d_i is non-zero for $3 \le i \le 6$. What happens if this is not the case?

Case 2: $v_1, w_2 \neq 0$, but $d_3 = d_4 = d_5 = d_6 = 0$. The system simplifies to

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 + d_2x_2y_2 = hl,$$

$$v_1x_1 + v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_2y_2 + w_3y_3 + w_4y_4 + w_5y_5 + w_6y_6 = l.$$

If any one of the coefficients v_3, \ldots, v_6 is non-zero, we can perform the same argument to conclude that $d_1 = 0$. A non-zero coefficient among w_3, \ldots, w_6 implies $d_2 = 0$. This would imply that hl = 0 and lead to the system

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

 $(\mathbf{v}^T \cdot \mathbf{x})(\mathbf{w}^T \cdot \mathbf{y}) = 0.$

It corresponds to the degenerate case (2) in Theorem 1.1. Therefore, we may assume (for example) that w_3, \ldots, w_6 are all zero.

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 + d_2x_2y_2 = hl,$$

$$v_1x_1 + v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_2y_2 = l.$$

Assume for now that $v_3 \neq 0$. If we set $x_2 = y_3 = 0$ by using Lemma 2.1, we end up with

$$x_1y_1 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 = hl,$$

$$v_1x_1 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_2y_2 = l.$$

The variables x_3 and y_2 have non-zero coefficients and appear only in linear equations. This allows us to use Lemma 2.7 to reduce the system to

$$x_1y_1 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 = hl,$$

where the number of solutions is bounded by $O(N^6(\log N)^2)$ by Lemma 2.6 as long as $d_1 \neq 0$.

The same argument works if one of v_4, \ldots, v_6 is non-zero. Therefore, we are doing fine, except when $d_1 = 0$ or $v_3 = v_4 = v_5 = v_6 = 0$.

Case 2.1:
$$v_3 = v_4 = v_5 = v_6 = 0$$
.

By replacing the auxiliary variables h and l, the system is now given by

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

 $d_1x_1y_1 + d_2x_2y_2 = (v_1x_1 + v_2x_2)(w_1y_1 + w_2y_2).$

Since the second equation is independent of the variables y_5 and y_6 , we can use Lemma 2.1 (iii) and Lemma 2.3 to bound the contribution of the first equation by $O(N^6 \log N)$ independent of the variables x_1, y_1, x_2, y_2 and consider the equation

$$d_1x_1y_1 + d_2x_2y_2 = (v_1x_1 + v_2x_2)(w_1y_1 + w_2y_2)$$

on its own. If the rank of the corresponding matrix is two, then Lemma 2.4 and Lemma 2.3 will give the correct upper bound. If, on the other hand, the rank is one, the we are again in the exceptional case (2) in Theorem 1.1.

Case 2.2:
$$d_1 = 0$$
 and $v_3 \neq 0$ (for example).

We took another small step forward in removing one more coefficient from the second bilinear equation. The system now looks like

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_2x_2y_2 = hl,$$

$$v_1x_1 + v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_2y_2 = l.$$

We can also assume that $d_2 \neq 0$ since the complementary case is covered earlier in 'Case 2'. The final case analysis is whether $w_1 = 0$ or not. If

 $w_1 = 0$, we obtain

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$(d_2x_2 - w_2h)y_2 = 0,$$

$$v_1x_1 + v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h.$$

If we insert the linear equation into the second equation, we see that the corresponding matrix has rank one. Therefore, we are in the exceptional case (2) of Theorem 1.1.

If $w_1 \neq 0$, on the other hand, we can set $x_1 = 0 = y_3$ by Lemma 2.1 and reduce the problem to

$$x_2y_2 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_2x_2y_2 = hl,$$

$$v_2x_2 + v_3x_3 + v_4x_4 + v_5x_5 + v_6x_6 = h,$$

$$w_1y_1 + w_2y_2 = l.$$

The variables x_3 and y_1 have non-zero coefficients and Lemma 2.7 allows us to remove the linear equations. The remaining system

$$x_2y_2 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_2x_2y_2 = hl,$$

has $O(N^6(\log N)^2)$ solutions by Lemma 2.6.

5. Parameter Case (II)

Now we have $B = D + \mathbf{v} \otimes \mathbf{e}_i + \mathbf{e}_i \otimes \mathbf{w}$. By a change of variables, we can assume that i = 1 and obtain the form

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_1x_1y_1 + d_2x_2y_2 + d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = y_1L(\mathbf{x}) + x_1M(\mathbf{y})$$

for the linear forms $L(\mathbf{x}) = \mathbf{v}^T \mathbf{x}$ and $M(\mathbf{y}) = \mathbf{w}^T \mathbf{y}$. The approach here is similar to the one in the previous section.

First we set $x_1 = y_1 = 0$ with the help of Lemma 2.1 and analyse the simpler system

$$x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_2x_2y_2 + d_3x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = 0.$$

If the d_i take on more than two values, we are done by Lemma 2.6. Otherwise, we can take linear combinations and simplify further to

$$x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_5x_5y_5 + d_6x_6y_6 = 0,$$

where $d_5 \in \{0, d_6\}$ (after a renaming of variables).

If $d_5 = d_6 \neq 0$, we are given the right upper bound by Lemma 2.6 again. Otherwise, we have $d_5 = 0$ and have found that our original system must

have the form

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$d_6x_6y_6 = y_1L(\mathbf{x}) + x_1M(\mathbf{y})$$

Here, we can use Lemma 2.1 to set $x_1 = 0$. This makes the second equation independent of y_2, \ldots, y_5 . By Lemma 2.1 (iii), this implies that we can simplify the system to

$$x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 = 0,$$

 $d_6x_6y_6 = y_1L(0, x_2, \dots, x_6).$

If $v_i \neq 0$ for some $i \in \{2, 3, 4, 5\}$, we can apply Lemma 2.1 (iii) again to remove the term $x_i y_i$ from the first equation and then change coordinates with Lemma 2.4 to obtain (here i = 2 for example)

$$x_3y_3 + x_4y_4 + x_5y_5 = 0,$$
$$d_6x_6y_6 = y_1x_2,$$

which has the right upper bound for the number of solutions by Lemma 2.6 as long as $d_6 \neq 0$.

This implies that we have the correct upper bound, except if $d_6 = 0$ or $L(\mathbf{x}) = L(x_1, 0, 0, 0, 0, x_6)$. A symmetric argument gives us the same conclusion with the condition $d_6 = 0$ or $M(\mathbf{y}) = M(y_1, y_6)$.

Case 1: $d_6 = 0$.

The system now simplifies to

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

 $y_1L(\mathbf{x}) + x_1M(\mathbf{y}) = 0.$

Here we need a slightly unusual procedure. We set $h = L(\mathbf{x})$ and $l = M(\mathbf{y})$ to lift the system to

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

 $y_1h + x_1l = 0,$
 $L(\mathbf{x}) = h,$
 $M(\mathbf{y}) = l.$

Now we perform Lemma 2.1 (iii) two times. One time with the set $\{x_2, \ldots, x_6\}$ and a second time with $\{y_2, \ldots, y_6\}$. This leaves us with the homogeneous system

$$x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$y_1h + x_1l = 0,$$

$$L(0, x_2, \dots, x_6) = 0,$$

$$M(0, y_2, \dots, y_6) = 0.$$

The second equation has four independent variables, which gives $O(N^2 \log N)$ by Lemma 2.3. The remaining system is an intersection of a diagonal bilinear form in $2 \cdot 5$ variables with two linear equations. The resulting bilinear

form has rank at least three and if those two equations aren't degenerate, we have the correct upper bound by Lemma 2.3.

Degenerate means here that one of the original linear forms L of M has to depend only on x_1 or y_1 . This would lead to a system of the shape

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

 $(v_1y_1 + M(\mathbf{y}))x_1 = 0,$

(or the equivalent for $M(\mathbf{y}) = w_1 y_1$), which has rank one in the second equation and corresponds to the exceptional case (2) in Theorem 1.1.

Case 2: $v_i = w_i = 0$ for $i \in \{2, 3, 4, 5\}$. In this case, we are left with the system

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$(v_1 + w_1)x_1y_1 + d_6x_6y_6 = v_6x_6y_1 + w_6x_1y_6.$$

We have seen this before in Case 2.1 of the previous section. Lemma 2.1 with Lemma 2.6 are sufficient to deal with it.

In this last case we have B = D + E, where (we can assume that) E has only non-zero entries in the upper left 3×3 corner. This corresponds to a system of the form

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$L_1(\mathbf{x}')y_1 + L_2(\mathbf{x}')y_2 + L_3(\mathbf{x}')y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = 0,$$

with $\mathbf{x}' = (x_1, x_2, x_3, 0, 0, 0).$

We use Lemma 2.1 to set $x_1 = x_2 = 0$ and $y_1 = y_2 = 0$, which reduces the problem to a diagonal one of the form

$$x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$l_{33}x_3y_3 + d_4x_4y_4 + d_5x_5y_5 + d_6x_6y_6 = 0.$$

By Lemma 2.6 we can deal with this situation, if the coefficients l_{33} , d_4 , d_5 , d_6 take on three different values. Otherwise, we can assume that $d_5 = d_6$.

Taking linear combinations in the original system, we therefore can simplify our problem to

$$x_1y_1 + x_2y_2 + x_3y_3 + x_4y_4 + x_5y_5 + x_6y_6 = 0,$$

$$L_1(\mathbf{x}')y_1 + L_2(\mathbf{x}')y_2 + L_3(\mathbf{x}')y_3 + d_4x_4y_4 = 0.$$

Since the second equations doesn't depend on y_5 and y_6 Lemma 2.1 (iii) reduces the problem further to

$$x_5y_5 + x_6y_6 = 0,$$

$$L_1(\mathbf{x}')y_1 + L_2(\mathbf{x}')y_2 + L_3(\mathbf{x}')y_3 + d_4x_4y_4 = 0.$$

Both equations are independent of each other. The first gives a bound of $O(N^2 \log N)$ by Lemma 2.3 and the second is fine as well by the same argument (with a coordinate change before), as long as the corresponding

matrix has rank at least two. Otherwise we are in the exceptional case (2) of Theorem 1.1.

7. Extension to s > 6

What happens, when the number of variables is larger than six? Either every linear combination of the two matrices has rank at most five, which brings us to the exceptional case (1) in Theorem 1.1, or we can find a change of coordinates, such that our system looks like

(7.1)
$$x_1 y_1 + \ldots + x_6 y_6 + d_7 x_7 y_7 + \ldots + d_s x_s y_s = 0,$$

$$L_1(\mathbf{x}) y_1 + \ldots + L_6(\mathbf{x}) y_6 + L_7(\mathbf{x}) y_7 + \ldots + L_s(\mathbf{x}) y_s = 0.$$

with $d_i \in \{0,1\}$. Setting $x_i = y_i = 0$ for all $i \geq 7$ by Lemma 2.1, we can use the result for s = 6 to see that we either get the general result or that we can add a multiple of the first equation to ensure that L_1, \ldots, L_6 are multiples of each other.

We can apply the same argument for any set of six variables for which $d_i \neq 0$. This results in the following structure for some value $f \geq 6$.

$$x_1 y_1 + \ldots + x_f y_f = 0,$$

 $L(\mathbf{x}) y_1 + \ldots + L(\mathbf{x}) y_f + L_{f+1}(\mathbf{x}) y_{f+1} + \ldots + L_s(\mathbf{x}) y_s = 0.$

By Lemma 2.1 this can be reduced to

$$x_1 y_1 + \ldots + x_f y_f = 0,$$

 $L_{f+1}(\mathbf{x}) y_{f+1} + \ldots + L_s(\mathbf{x}) y_s = 0,$

and we are done, as long as these linear forms L_j are not all multiples of each other.

In the complementary case, the rank of the second equation in (7.1) is at most two. If it is less than two, we are done. Otherwise we perform a suitable change of coordinates, swap the equations, and obtain the form

$$x_1y_1 + x_2y_2 = 0,$$

$$\tilde{L}_1(\mathbf{x})y_1 + \ldots + \tilde{L}_s(\mathbf{x})y_s = 0.$$

for some other linear forms $\tilde{L}_1, \ldots, \tilde{L}_s$. Since the second equation must have rank at least four, we can find two linear forms L_i and L_j for i > j > 2, which are linearly independent. An application of Lemma 2.1 gives the system

$$x_1y_1 + x_2y_2 = 0,$$

$$L_i(\mathbf{x})y_i + L_i(\mathbf{x})y_i = 0$$

and we are done. This is the end of the proof for Theorem 1.1.

APPENDIX A. PROOF OF LEMMA 3.1

Let B be a matrix with off-rank one. By permuting variables, if necessary, we can assume that B has the form

$$B = \begin{pmatrix} a & r & \mathbf{m}^T \\ s & b & \mathbf{n}^T \\ \mathbf{v} & \mathbf{w} & C \end{pmatrix},$$

where $r \neq 0$ and $\mathbf{m}, \mathbf{n}, \mathbf{v}, \mathbf{w} \in \mathbb{Z}^{s-2}$. The following lemma is the first step to understand the structure of B.

Lemma A.1. For the above matrix we have

$$C = \mathbf{wm}^T / r + D$$
,

where D is a diagonal matrix.

Proof. Consider the 2×2 submatrix $\begin{pmatrix} r & m_j \\ w_i & c_{ij} \end{pmatrix}$ for some $i \neq j$. Since the off-rank is one, this matrix has rank at most one. Since $r \neq 0$ it must be at least one. A short calculation shows that $c_{ij} = w_i m_j / r$.

Since the off-rank of B is one, we also get that $\mathbf{v} = \lambda \mathbf{u}, \mathbf{w} = \mu \mathbf{u}$ and $\mathbf{m} = \alpha \mathbf{k}, \mathbf{n} = \beta \mathbf{k}$ for some $\mathbf{u}, \mathbf{k} \in \mathbb{Z}^{s-2} \setminus \{\mathbf{0}\}$ and $\lambda, \mu, \alpha, \beta \in \mathbb{Q}$. We obtain

$$B = \begin{pmatrix} a & r & \alpha \mathbf{k}^T \\ s & b & \beta \mathbf{k}^T \\ \lambda \mathbf{u} & \mu \mathbf{u} & C \end{pmatrix},$$

Case 1: $\alpha = 0$.

This implies that C is diagonal.

Case 1.1: There are $i \neq j$ such that $u_i \neq 0$ and $k_j \neq 0$

Consider the off-diagonal matrix $\begin{pmatrix} s & \beta k_j \\ \lambda u_i & c_{ij} \end{pmatrix}$. Since $c_{ij} = 0$, we conclude that $\beta = 0$ or $\lambda = 0$.

Case 1.1.1: $\lambda = 0$.

We are in case (ii) of Lemma 3.1 and done.

Case 1.1.2: $\beta = 0$ and $\lambda \neq 0$.

If $\mu = 0$, we are in case (ii) again. Otherwise we have to show that we can choose entries x, y in

$$\begin{pmatrix} x & r & \mathbf{0}^T \\ s & y & \mathbf{0}^T \\ \lambda \mathbf{u} & \mu \mathbf{u} & O \end{pmatrix},$$

such that the resulting matrix has rank one. Choose $x = r\lambda/\mu$ and $y = s\mu/\lambda$. It follows that $B = D + \mathbf{v} \otimes \mathbf{w}$ for some $\mathbf{v}, \mathbf{w} \in \mathbb{Q}^s$ and a diagonal matrix D, which corresponds to case (i).

Case 1.2: For all $i \neq j$ we have $u_i = 0$ or $k_j = 0$.

The condition implies that $\mathbf{u} = \mathbf{0}, \mathbf{k} = \mathbf{0}$ or that there is at most one index i such that $u_i \neq 0$ and $k_i \neq 0$.

Case 1.2.1: u = 0.

We are in case (ii) of Lemma 3.1.

Case 1.2.2: k = 0.

This brings us back to the Cases 1.1.1 and 1.1.2.

Case 1.2.3: $u_j = k_j = 0$ for all $j \neq i$ for some fixed i.

This implies that only the 1st, 2nd and ith row/column have non-zero non-diagonal entries, which brings us into case (iii).

Case 2: $\mu = 0$.

This is completely analogous to Case 1.

Case 3: $\alpha \neq 0$ and $\mu \neq 0$.

Now there is at least one non-diagonal entry $c_{ij} \neq 0$. Consider the matrix $\begin{pmatrix} s & \beta k_j \\ \lambda u_i & c_{ij} \end{pmatrix}$. We know that $c_{ij} = \mu u_i \alpha k_j / r$. This implies that the matrix

can have rank one only if $s = \lambda \beta u_i k_j / c_{ij} = \frac{r \lambda \beta}{\mu \alpha}$. If we consider B modulo diagonal matrices, we see that we can choose x and y such that

$$\begin{pmatrix} x & r & \alpha \mathbf{k}^T \\ r\lambda \beta/\mu \alpha & y & \beta \mathbf{k}^T \\ \lambda \mathbf{u} & \mu \mathbf{u} & \alpha \mu \mathbf{u} \mathbf{k}^T/r \end{pmatrix},$$

has rank one by setting $x = r\lambda/\mu$ and $y = r\beta/\alpha$. This gives us case (i) in Lemma 3.1.

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